## CLAIMS

## What is claimed is:

1	1. A	method	of locating an entry in a forwarding database corresponding to a longest	
2	match of an address, the method comprising the steps of:			
3		a)	applying a mask to the address to determine a masked address that is to be	
4 ·			used for purposes of locating a matching entry in the forwarding database;	
5		b)	searching the forwarding database for an entry that matches the masked	
6			address;	
7		c)	performing an address-sensitive decimation of the mask to produce a new	
8			mask; and	
9		d)	until a predetermined condition has been met, repeating steps a-c with the	
10			new mask.	
1	2.	The n	nethod of claim 1, further comprising the step of performing a hash function	
2	4.		e masked address to produce an index into a hash table.	
1	3.	The r	nethod of claim 2, wherein the step of performing an address-sensitive	
2		decimation of the mask to update the mask comprises shortening the mask such that		
3		the a	pplication of the mask to the address will exclude enough data to cause the	
4		hash	function to produce a result that is different from the former index.	
1	4.	The	method of claim 3, wherein the step of shortening the mask such that the	
2			cation of the mask to the address will exclude enough data to cause the hash	
3		func	tion to produce a result that is different from the former index comprises	
4		shor	tening the length of the mask to X-Y bits, where X represents the bit position	

5		of the most significant bit in the address, and where Y represents the bit position of
6		the least significant bit containing a one.
1	5.	The method of claim 1, wherein the predetermined condition comprises a matching
2		entry has been located or the length of the mask is equal to a predetermined
3		threshold.
1	6.	The method of claim 1, wherein one or more new masks are generated in parallel
2		with the step of searching the forwarding database for a matching entry.
1	7.	The method of claim 1, wherein each entry in the forwarding database includes N-
2		bits of address and an associated mask length encoded in N+1 bits of address
3		information, and the step of searching the forwarding database for a matching entry
4		comprises the steps of:
5		determining the mask length associated with an entry based upon the location of a
6		mask length identification marker within the address information;
7		comparing the mask length associated with the entry with the mask that has been
8		applied to the address; and
9		if the mask length associated with the entry is less than or equal to the mask that has
10		been applied to the address, then comparing a significant portion of the
11		address information with the masked address, the significant portion of the
12		address information being identified by the mask length associated with the
13		entry.
1	8.	The method of claim 7, wherein the mask length identification marker comprises the
2		least significant bit in the address information containing a one, and wherein the
2		step of determining the mask length associated with an entry based upon the

	4	locati	on or a	mask length identification marker within the address information				
	5	comprises scanning the address information to locate the bit position of the least						
	6	signi	ficant bi	t containing a one.				
	1	9. A method	d of loca	ating an entry in a forwarding database corresponding to a longest				
	2	match of a search key, the method comprising the steps of:						
	3	a)	searc	hing the forwarding database for an entry that matches the search key;				
	4		and					
	5	b)	if no	entry matches the search key, then				
	6		1)	scanning the search key to locate the least significant bit containing a				
	7			one,				
	8		2)	shortening the search key to exclude the least significant bit				
	9			containing a one,				
	10		3)	searching the forwarding database for an entry that matches the				
	11			search key, and				
	12		4)	repeating steps 1-3 until the search key is equal to a predetermined				
	13		•	length or until the longest match is located.				
	1	10. The	method	of claim 9, wherein the step of shortening the search key to exclude the				
	2	least significant bit containing a one comprises the steps of:						
	3	providing a mask of length X-Y, where X represents the bit position of the most						
	4		sign	ificant bit in the search key, and where Y represents the bit position of				
	5		the l	east significant bit containing a one; and				
	6	ann	lving the	e mask to the search key.				

i	11.	The method of claim 9, wherein the search key comprises a destination Internet
2		Protocol (IP) address.
1	12.	The method of claim 19, wherein the search key comprises a source Internet
2		Protocol (IP) address.
1	13.	A method of forwarding data comprising the steps of:
2		receiving data at a port;
3		extracting an address from the data;
4		searching a forwarding database for a longest match for the address by
5		comparing a portion of the address indicated by a mask to entries in the
6		forwarding database, and
7		progressively shortening the mask based upon the address until a matching
8		entry is located; and
9		forwarding the data to a destination associated with the matching entry.
1	14.	The method of claim 13, further comprising the step of performing a hash function
2		on the portion of the address to produce an index into a hash table.
1	15.	The method of claim 14, wherein the step of progressively shortening the mask
2		based upon the address comprises, for each search iteration, shortening the mask
3		such that the application of the mask to the address will exclude enough data to
4		cause the hash function to produce a result that is different from the former index.
1	16.	The method of claim 13, wherein the step of progressively shortening the mask
2		based upon the address until a matching entry is located comprises, for each search
3		iteration, shortening the length of the mask to X-Y bits, where X represents the bi

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4		position of the most significant bit in the address, and where Y represents the bit
5		position of the least significant bit containing a one.
1	17.	The method of claim 13, wherein each entry in the forwarding database includes N-
2		bits of address and an associated mask length encoded in N+1 bits of address
3		information, and the step of comparing a portion of the address indicated by a mask
4		to entries in the forwarding database comprises the steps of:
5		determining the mask length associated with an entry based upon the location of a
6		mask length identification marker within the address information;
7		comparing the mask length associated with the entry with the mask; and
8		if the mask length associated with the entry is less than or equal to the mask, then
9		comparing a significant portion of the address information with the portion
10		of the address indicated by the mask, the significant portion of the address
11		information being identified by the mask length associated with the entry.
1	18.	The method of claim 17, wherein the mask length identification marker comprises
2		the least significant bit in the address information containing a one, and wherein the
3		step of determining the mask length associated with an entry based upon the
4		location of a mask length identification marker within the address information
5		comprises scanning the address information to locate the bit position of the least
6		significant bit containing a one.
1	19.	The method of claim 13, wherein the address comprises a 32-bit or a 128-bit
2		Internet Protocol (IP) address.

match of a search key, the method comprising the steps of:

A method of locating an entry in a forwarding database corresponding to a longest

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3		a)	perfo	ming a hash function on the search key to produce a current index into		
4			a hasł	table;		
5		b)	searcl	ning a first bin in the hash table identified by the current index for an		
6			entry	that matches the search key; and		
7		c)	while	no entry is found that matches the search key and while the search key		
8			is gre	ater than a predetermined length, with each subsequent search iteration		
9			perfo	rming the steps of		
10			1)	shortening the search key to exclude just enough data to cause the		
11				hash function to produce a result that is different than the current		
12				index;		
13			2)	updating the current index with the result of the hash function on the		
14				shortened search key; and		
15			3)	searching a different bin in the hash table that is identified by the		
16				current index.		
1	21.	A m	ethod o	f locating an entry in a forwarding database corresponding to a longest		
2		mate	ch of a s	h of a search key, the method comprising the steps of:		
3		a)	generating indices for a hash table performing a hash function on the search			
4			key	to produce a current index into a hash table;		
5		b)	sear	ching a first bin in the hash table identified by the current index for an		
6			entr	y that matches the search key; and		
7		c)	whi	le no entry has been found that matches the search key and while the		
8			leng	th of the search key is greater than a predetermined length, with each		
9			sub	sequent search iteration performing the steps of		

	10		1)	shortening the search key to exclude just enough data to cause the
	11			hash function to produce a result that is different than the current
	12			index;
	13		2)	updating the current index with the result of the hash function on the
	14			shortened search key; and
	15		3)	searching a different bin in the hash table that is identified by the
	16			current index.
	1	22.	A networking	g device comprising:
	2		a backplane;	and
	3		a plurality of	input/output (I/O) interfaces coupled to the backplane,
	4		each	of the plurality of I/O cards comprising
	5		a plu	rality of ports,
	6		a for	warding and filtering mechanism coupled to the plurality of ports, the
	7			forwarding and filtering mechanism configured to forward data
	8			based upon the results of a longest match search of a forwarding
è	9			database for an entry corresponding to an address contained within
	10			the data, where:
	11			a mask is applied to the address to determine a masked address to be
	12			used for purposes of searching the forwarding database,
	13			the forwarding database is searched for entries that match the
	14			masked address, and
	15			subsequent masks are produced based upon an address-sensitive
	16			decimation of the mask.